

PFA(S) implies there are many S-names

Alan Dow

Department of Mathematics and Statistics
University of North Carolina Charlotte

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what I mostly want to talk about today is our result

MM(S) implies that forcing with S makes all normal locally
compact spaces \aleph_1 -CWH

because the proof is so different

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from this day forward (and since a few years back)

S refers to a fixed Souslin tree in $L[A]$ that is **coherent**:

S is a subtree of $\omega^{<\omega_1}$ satisfying that, for each $\alpha \in \omega_1$ and $s \in S_\alpha$
and $t \in \omega^\alpha$, $t \in S$ iff $\{\beta < \alpha : t(\beta) \neq s(\beta)\}$ is finite

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Definition

A poset \mathbb{P} is S-preserving means that S is still a Souslin tree in the forcing extension by \mathbb{P} , and, as usual

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$MA_{\mathcal{C}}(\omega_1)$ for a class \mathcal{C} of posets means for $\mathbb{P} \in \mathcal{C}$: for every \aleph_1 many dense sets, there is a (generic) filter meeting them all.

Souslin preserving method

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where SA_{ω_1} is formulated (we could say $MA(S)$):

S is Souslin and $MA_{\mathcal{C}}(\omega_1)$ holds where \mathcal{C} is all S -preserving ccc posets

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also $MM(S)$ is investigated in Miyamoto's JSL paper

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\exists LCN non \aleph_1 -CWH LCN implies \aleph_1 -CWH

basics of PFA(S)

If \mathbb{P} is proper and $p \in \mathbb{P}$ is (M, \mathbb{P}) -generic, then we want to prove that $S \Vdash p$ is $(M[g], \check{\mathbb{P}})$ -generic – *same(?)* as $\mathbb{P} \times S$ is proper because then we can apply PFA(S) to \mathbb{P} ; to get a useful S -name.
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Lemma

If $D \in M$ is a dense subset of $\mathbb{P} \times S$, then

$$\dot{E} = \{r \in \mathbb{P} : (\exists s \in g) (r, s) \in D\}$$

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Let \mathbb{P} be proper and $p \in \mathbb{P}$ be (M, \mathbb{P}) -generic. Then coherence solves one problem for us, but not another.

Lemma

If $s_1, s_2 \in M \cap S$ and if s_1 forces that p is $(M[g], \check{\mathbb{P}})$ -generic, then so does s_2 .

this is because of (there's only one generic extension)

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If $\dot{A} \in M$ is any S -name of a subset of \mathbb{P} , then there is another S -name $\dot{B} \in M$ such that if s_2 is in any S -generic g_2 , there is an S -generic $g_1 \ni s_1$ such that $\dot{A}[g_2] = \dot{B}[g_1]$. use \dot{A}_{s_1, s_2} to denote \dot{B}

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Proof.

by extending, we assume that s_1 and s_2 are in same S_β .

Simply $\dot{B} = \{(p, s_1 \oplus t) : (p, t) \in \dot{A}, s_2 \not\subseteq t\} ; s_1 \oplus t = s_1 \cup (t \setminus s_2)$
and define $g_1 = \{s_1 \oplus t : s_2 \subset t \in g_2\}$. □

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of course $\text{PFA}(S)$ implies there is an S -space, but T. Yorioka proved that none remain an S -space in $\text{PFA}(S)[g]$

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$\{p_\alpha : \alpha \in \omega_1\}$ unravels as $\dot{Y} = \{(x_\delta, s_\delta) : \delta \in C\}$, a name such that

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all is fine if we can prove $\mathbb{P} \times S$ is proper

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s_{M_1} forces a value on $\dot{W}_{x_{M_1}} \cap M_1$ – call it $\dot{W}_{x_{M_1}}[s_1]$

all is fine if we can prove $\mathbb{P} \times S$ is proper

consider e.g. $(p, s_0) \in D \in M \prec H(\theta)$ and

$p = \{ (M_1, (x_1, s_1)), (M_2, (x_2, s_2)), (M_3, (x_3, s_3)) \}$ where
 $M_1 = M \cap H(\kappa)$

and as in picture for S-positions

the PFA no S-space picture

$\{x_1, x_2, x_3\}$ from a tree T

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$\{x_1, x_2, x_3\}$ from a tree T
separated by models
means T is ω_1 -branching

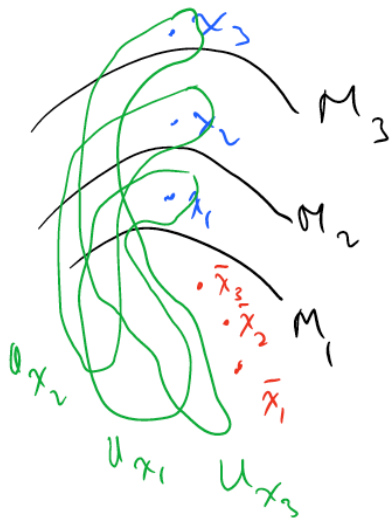
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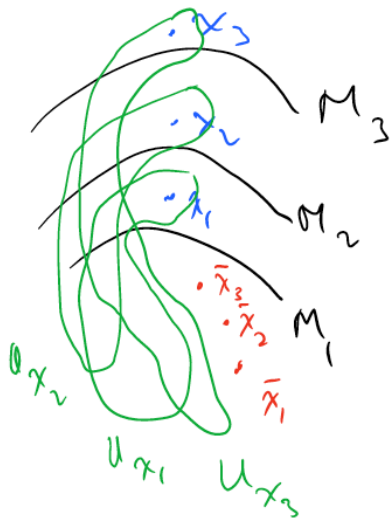
$\{x_1, x_2, x_3\}$ from a tree T
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need (M_1, \mathbb{P}) -generic

we must reflect into M_1



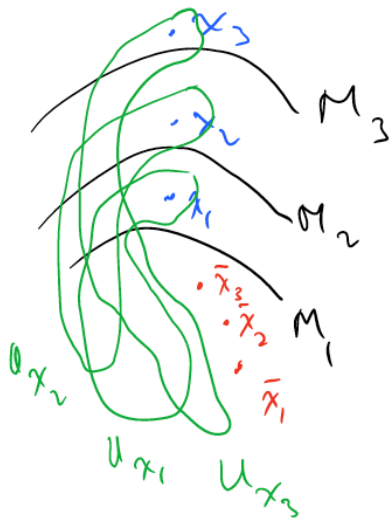
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using that $U_{x_1} \cup U_{x_2} \cup U_{x_3}$
 has countable closure

and that $M_1 \cap T$ does not
 we can find $\{\bar{x}_1\} \in M_1 \cap T$



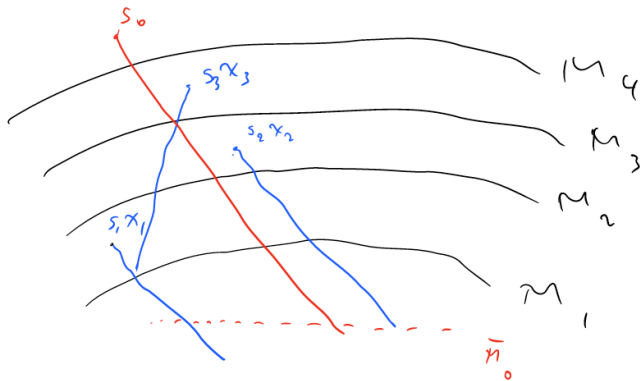
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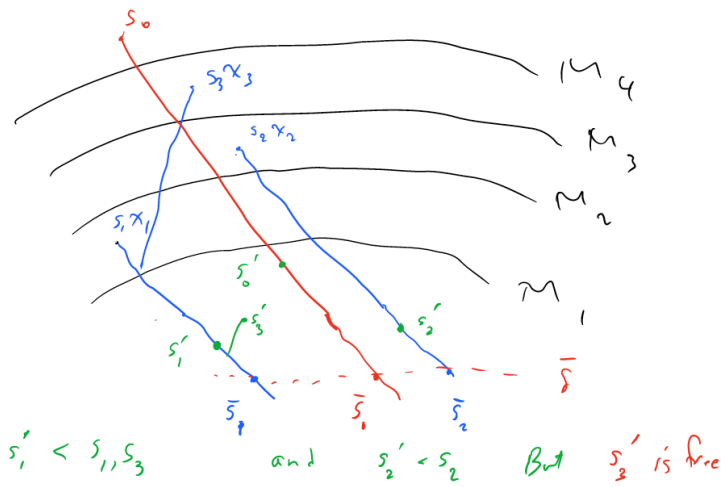
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 we can find $\{\bar{x}_1\} \in M_1 \cap T$
 with $\bar{x}_1 \notin U_{x_1} \cup U_{x_2} \cup U_{x_3}$

$\{x : \{\bar{x}_1, x\} \in T\} \in M_1$
 repeat to find \bar{x}_2

now picture $s_0 \in S$ and $S \times X$ pairs from \mathbb{P}



s_0 is saying “you must work in my universe” but \mathbb{P} has conditions as we again try to reflect into M_1



we are forced to have $s'_0 < s_0$ which, by elementarity, then determines entire structure; discuss W_x 's next

We set $\dot{E} = \{r : (\exists s_r) (r, s_r) \in D\}$ and (by $M \prec H(\theta)$)
have that $s_0 \Vdash \dot{A} = \{x_1^r : r \in \dot{E}\}$ is uncountable and in M

and this is fine (enough) for finding $x_1^r \in M \setminus W_{x_1}$ such that
 $s_0 \Vdash x_1^r \in \dot{A}$ we can even get $x_1^r \notin W_{x_3}$

and $s_2 \Vdash (\exists r \in \dot{E} \cap M) x_2^r \notin W_{x_2}$ is (fairly) standard

but for \dot{W}_{x_1} , we must (and can) use that $s_1 \Vdash \dot{A}_{\bar{s}_0, \bar{s}_1}$ is uncountable

Unsolved is how to get required x_1^r not in $\dot{W}_{x_1}[s_1] \cup \dot{W}_{x_3}[s_3]$
because this may even cover $M \cap X$.

Moore-Mrowka under PFA(S)?

Remark

PFA(S) implies there is a compact sequential X that, after forcing with S , has uncountable tightness; **so we can't simply apply Moore-Mrowka in PFA(S)[g] model**

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PFA method for Moore-Mrowka

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thus compact countably tight X has cardinality at most c

Lemma (1. first reduction)

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We now assume that ω_1 is a closed discrete subset of a locally compact space X of weight \aleph_1 . We have $\{U(\alpha, \xi) : \xi \in \omega_1\}$ a neighborhood base (with compact closures) at $\alpha \in \omega_1$. For convenience $\overline{U(\alpha, \xi + 1)} \subset U(\alpha, \xi) \subset U(\alpha, 0)$.

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Proposition [after forcing with S]

There is a cub C_0 so that for all $\delta \in C_0$, there is a $\beta(\delta) < \delta_C^+$ such that $\omega_1 \cap \overline{\bigcup_{\alpha < \delta} U(\alpha, 0)} \subset \beta(\delta)$.

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Proof.

This is because $\overline{\bigcup_{\alpha < \delta} U(\alpha, 0)}$ has a dense Lindelof subspace and so by Lemma 2, its closure has countable extent. \square

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Proof.

Larson-Tall prove that in a forcing extension by S , a closed discrete set of \aleph_1 many points of countable character is separated if it is normalized. If A_C is not stationary, we can shrink to C_1 so that the quotient space obtained by collapsing each $Z(\alpha, C_1)$ to a point will result in the image of ω_1 being closed discrete and normalized. \square

Now we use the generic g viewed as $\in \omega^{\omega_1}$

Definition

In the ground model (of $\text{MM}(S)$), let $\{C_\gamma : \gamma \in \omega_2\}$ be a base for the cub filter, chosen so that, for all $\zeta < \gamma$, $C_\gamma \setminus C'_\zeta$ is countable.

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now fix S-names for everything

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Lemma

$\{E_\gamma : \gamma \in \omega_2\}$ is a family of stationary subsets of ω_1 .

Proof.



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- 5 there is a single $\bar{s} \in S$ so that $\{s_\gamma : \gamma \in I\}$ is dense above \bar{s} .



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- ⑩ there is an L such that $s_\gamma \upharpoonright \delta_{C_\gamma}^+$ forces that $\dot{U}(\bar{\beta}, 0)$ is disjoint from $\bigcup \{ \dot{Z}(\alpha, C_{\zeta(\gamma)}) : \alpha < \delta \text{ and } \sigma_{C_\gamma}(\alpha) \geq L \}$ because $\alpha_{C_{\zeta(\gamma)}}^+ < \delta$ for all $\alpha < \delta$ and $\dot{f}_{C_\gamma}(\dot{U}(\bar{\beta}, 0))$ is bounded.



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and, since s_{γ_L} forces that $f_{C_\gamma}(\bar{\beta}) = L$, it forces

$$\bar{\beta} \text{ is not in } \overline{\bigcup \{ \dot{Z}(\alpha, C_{\zeta(\gamma_L)}) : \alpha \in \delta, \sigma_{C_\gamma}(\alpha) < L \}}.$$

