

Turning ternary relations into antisymmetric betweenness relations

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The concept of betweenness

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- Natural to regard such a relation as a ternary predicate $[a, b, c]$, where $(a, b, c) \in X^3$.
- Birkhoff (1948) defined the betweenness relation $[\cdot, \cdot, \cdot]_o$ on a partially ordered set (X, \leq) as an extension of that given above.

Examples of betweenness: partial orders

Definition

In a partially ordered set (X, \leq) with $d \leq e \in X$, define the order interval $[d, e]_o = \{x \in X : d \leq x \leq e\}$.

- If each pair of elements in X has a common lower bound and a common upper bound in X , then say that $[a, b, c]_o$ if b belongs to *each* order interval that also contains a and c .

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- Metric space (X, d) (1928):
define $[a, c, b]_M$ if $d(a, c) + d(c, b) = d(a, b)$.
- Natural alliance between intervals $[a, b]$ and ternary predicates $[a, c, b]$, in that we intend $c \in [a, b]$ iff $[a, c, b]$.

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while the largest is $X_{\top} := X^3 \setminus \{[a, b, a] \mid a \neq b\}$.

Bankston's insight: road systems

Definition

A road system on a nonempty set X is a family \mathcal{R} of nonempty subsets (*roads*) of X such that

- (i) $\{a\} \in \mathcal{R}$ for all $a \in X$,
- (ii) for all $a, b \in X$, there is $R \in \mathcal{R}$ such that $a, b \in R$.

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Define $[a, c]_{\mathcal{R}} = \bigcap \{R \in \mathcal{R} : a, c \in R\}$.

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A road system characterization of betweenness

Theorem (Bankston, 2011)

A ternary relation $[\cdot, \cdot, \cdot]$ on a set X can be generated from a road system if and only if $[\cdot, \cdot, \cdot]$ is an R -relation.

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Antisymmetric R-relations

A road system \mathcal{R} is *separative* if for any $a, b, c \in X$ with $b \neq c$, there is some $R \in \mathcal{R}$ such that either $a, b \in R$ and $c \notin R$ or $a, c \in R$ and $b \notin R$.

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A category \mathbf{T} of ternary relations

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monotone functions: for objects $(X, [\cdot, \cdot, \cdot]_X)$ and $(Y, [\cdot, \cdot, \cdot]_Y)$ then $f : X \rightarrow Y$ is a morphism provided $[a, b, c]_X \Rightarrow [f(a), f(b), f(c)]_Y$.

Some notation and definitions

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The inclusion functor $\mathbf{R}_1 \hookrightarrow \mathbf{T}$

(R1) Reflexivity: $[a, b, b]$

The left adjoint is given by $(X, [\cdot, \cdot, \cdot]) \mapsto (X, [\cdot, \cdot, \cdot]')$ where

$$[\cdot, \cdot, \cdot]' = [\cdot, \cdot, \cdot] \cup \{[a, b, b] \in [\cdot, \cdot, \cdot] \mid a, b \in X\}.$$

Denote by L_1 .

The inclusion functor $\mathbf{R}_2 \hookrightarrow \mathbf{T}$

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The left adjoint exists - and is more involved. Call it L_3 .

The inclusion functor $\mathbf{R}_4 \hookrightarrow \mathbf{T}$

(R4) Transitivity: $[a, b, c] \wedge [a, d, c] \wedge [b, x, d] \Rightarrow [a, x, c]$

Has a left adjoint - call it L_4 .

Adjoint operators

Notice that the compositions $L_1 \circ L_2$ and $L_2 \circ L_1$ are not the same. The operator L_2 (closure under symmetry) does not preserve R1 (reflexivity).

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A less trivial example is given by L_3 and L_4 .

In fact, $L_4 \circ L_3 \circ L_1 \circ L_2$ defines the left adjoint to $\mathbf{R} \leftrightarrow \mathbf{T}$.

The subcategory **A** of antisymmetric R-relations

Antisymmetry: $[a, b, c] \wedge [a, c, b] \implies b = c$.

Question: does the inclusion functor $\mathbf{A} \hookrightarrow \mathbf{R}$ have a left adjoint?

Yes - demanding a change of underlying set; call it L_A .

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... but not necessarily R4.

And L_4 may not preserve antisymmetry.

Theorem

The left adjoint is the direct limit of applying L_4 after L_A ω -many times.

Mar fhocal scoir

Given a lattice (X, \leq) , define $[a, b]_L = \{x : a \wedge b \leq x \leq a \vee b\}$.

Lemma

Let $(X, [\cdot, \cdot, \cdot])$ be the R -relation generated from the lattice intervals (roads) described above.

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Lemma

Let $(X, [\cdot, \cdot, \cdot])$ be the R -relation generated from the lattice intervals (roads) described above.

Then (X, \leq) is distributive if and only if

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Lemma

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Then (X, \leq) is distributive if and only if $(X, [\cdot, \cdot, \cdot])$ is antisymmetric.